



A Framework accommodating Categorized Multiprocessor Real-time Scheduling in the RTSJ

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1. Backgrounds

- RTSJ and its implementations provide very good environment for real-time Java applications
 - User-implemented schedulers can be applied to make scheduling decisions
 - However, it does not have any multiprocessor related features
 "How can we make it useful for multiprocessor based schedulers?"
- Scheduling on multiprocessor systems is far more difficult than on single processor systems
 - There is no the Almighty solution for this problem yet
- Several operating systems are capable of multiple schedulers within their kernels





1.1 Categorization of Scheduling Algorithms

- In 2004, Carpenter et al.[1] presented a taxonomy for multiprocessor scheduling algorithms in two-dimensional space
 - The complexity of the priority scheme
 - Static
 - Dynamic, but fixed within a job
 - Fully dynamic
 - The degree of migration allowance
 - No migration
 - Migration only at job boundaries
 - Unrestricted migration

unrestricted

par	itioned/static			fully dynar	mic
migr		1: Static	2: Job-level dynamic	3: Fully dynamic	
ation degree	1: Partitioned	(1,1)-restricted	(2,1)-restricted	(3,1)-restricted	
	2: Restricted migration	(1,2)-restricted	(2,2)-restricted	(3,2)-restricted	
	3: Unrestricted migration	(1,3)-restricted	(2,3)-restricted	(3,3)-restricted	





1.1 Categorization of Scheduling Algorithms

- The progenitors of all scheduling algorithms[2] (grouped by priority dynamics)
 - Static: RM
 - Job-level dynamic: EDF
 - Fully dynamic: LLF
- The basic schedulers for each category

gree	Global RM	Global EDF			
ation de	Restricted migration RM	Restricted migration EDF	Restricted migration LLF		
migra	Partitioned RM	Partitioned EDF	Partitioned LLF		

priority complexity





1.2 Scheduler Frameworks

- Embracing multiple scheduling algorithms
 - Operating system
 - Kernel
 - Linux
 - LITMUSRT
 - ChronOS
 - Kernel/Application
 - RESCH
 - Middleware framework
 - RTSJ implementations
 - JEOPARD





1.3 RTSJ

- The Real-Time Specification for Java
 - Current official spec: 1.0.2
 - Real-time thread scheduling
 - Memory management
 - Resource sharing
 - Asynchronous event handling
 - Recent 'Alpha' release: 1.1 Alpha 6
 - Processor affinity for threads and event handlers
 - Minor updates, improvements and bugfixes





1.3 RTSJ

- Implementation of the RTSJ e.g) Java Real-Time System (v2.2u1)
 - A real-time virtual machine
 - Java HotSpot Server VM + @ (RT features)
 - RTSJ class libraries
 - RTSJ 1.0.2 (full)
 - Java class libraries
 - Java SE 5





1.3 RTSJ

• Available RTSJ implementations

Implementation	RTSJ Ver.	Vendor	Remarks
RTSJ RI	1.1a6 / 1.0.2(6)		Last released in JAN 2009
Sun Java RTS	1.0.2	Oracle	
IBM WebSphere RT	1.0.2	IBM	
JEOPARD	1.1 (JSR282)	JEOPARD Consortium	Uses parallel JamaicaVM from aicas GmbH
Perc Pico	JSR302	aonix	J2ME environment
OVM	1.0.1	Purdue University	Under BSD License





2. The CMRF

- The Categorized Multiprocessor Real-time scheduling-supporting middleware Framework
 - Goal: To provide functions to embrace the progenitor algorithms of each category
 - With minimum changes on existing system
 - Environment
 - OS: Linux, kernel 2.6.x
 - Native POSIX Thread Library
 - (PREEMPT_RT or equivalent kernel preemption patches)
 - RTSJ version: 1.0.2 (partial implementation)
 - RealtimeThread
 - Time and timers
 - Scheduler and scheduling parameters





2.1 The CMRF - Structures

- SCHED_FIFO base scheduler
 - Provided by Linux kernel
 - The first-come thread with the highest priority is served first.
 - Serving PriorityScheduler in the RTSJ
- Scheduling 'Schedulables' of the RTSJ
 - Only RealtimeThread is currently supported.
 - Direct use of SCHED_FIFO for thread scheduling
 - sched_setparam: Priority changes
 - sched_setaffinity: Thread migration
- Timers as sleep functions
 - The basic timer is implemented using nanosleep() function provided by underlying OS via Java Native Interface



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2.1 The CMRF - Structures

- The structure blocks of the CMRF
 - Native block
 - Java Runtime Environment
 - Framework and RTSJ block
 - Application schedulers
- Scheduling sequence
 - Scheduling decisions are made by application schedulers
 - Within the decision making, the eligible thread is dispatched through the framework
 - The dispatched thread is scheduled by SCHED_FIFO







2.1 The CMRF - Structures

- Scheduling events
 - Arrival of a new RealtimeThread
 - Release of next period of a RealtimeThread
 - Finishing a job
 - Timer







- Evaluation
 - Task parameters
 - Cost: ≤ 10ms
 - Deadline: 100ms
 - Period: 100ms
 - Total jobs: 100,000 releases / processor
 - Test system
 - Two Intel Xeon E5506 2.13GHz processors, 4 cores per processor
 - 8 total processors
 - Four AMD Opteron 6176 SE 2.3GHz processors, 12 cores per processor
 - 48 total processors
 - Ubuntu Linux 10.04.4, kernel 2.6.31
 - With PREEMPT_RT patch applied
 - OpenJDK6
 - HotSpot Server VM





- Deadline-miss ratio test, partitioned
 - RM starts to miss deadlines from utilization of 7.2/33.6 (90%/70%) with miss ratio of 0.95% and 0.31%
 - EDF and LLF starts missing from 7.2/43.2 (90%) with ratio of 0.39%/0.09%(EDF), 0.31%/0.05%(LLF) respectively







- Deadline-miss ratio test, restricted migration
 - RM starts to miss deadlines from utilization of 6.4 (80%), ratio of 0.5%
 - EDF and LLF starts missing from 6.4/43.2 (80%/90%), 6.4/48(80%/100%) with ratio of about 0.1% and 0.2%







- Deadline-miss ratio test, unrestricted migration
 - RM starts to miss deadlines from utilization of 6.4/38.4 (80%) with ratio of about 0.7%
 - EDF starts missing at 8/48 (100%), ratio of 0.83%/0.98%







- Elapsed time of each algorithms
 - Measurement of time taken for each algorithm
 - Partitioned schemes takes from 81~152us
 - Dispatch() call takes 90~100us, however, partitioned algorithms do not use affinity assign features in the call
 - As calling Dispatch() frequently, the algorithm gets slower
 - Restricted migration schemes takes around 454~628us
 - Global migration schemes takes 643~801ms







- Jitter test (within single processor)
 - Measures release overhead
 - Jitter time = [targeted release time actual release time]

– Parameters

- Partitioned
- Utilization: 0.2
- Total releases: 160,000 jobs / processor

Intel Xeon E550	6	(in	n nanosec.)	AMD Opteron 6	176SE	(i	n nanosec.)
	JRTS	OVM	CMRF		JRTS	OVM	CMRF
Jitter, Min.	2,438	2,529	2,823	Jitter, Min.	1,457	58	1,308
Jitter, Max.	144,426	190,048	52,404	Jitter, Max.	2,295,608	7,672,276	2,220,614
Average jitter	53,546	62,201	52,404	Average jitter	302,836	417,186	293,807
Deviation	20,286	23,017	14,053	Deviation	180,192	359,516	117,969





3. Conclusion

- Two-dimensional categorization needs:
 - Priority axis: priority parameters
 - Migration axis: affinity setting
- This can be done by using corresponding system calls and SCHED_FIFO policy
 - Thread scheduling using FIFO policy replaces Schedulables scheduling on the RTSJ
- No other specially built JVM is needed to schedule real-time Java threads on a system





Appendix – Issues and Discussions

- Other major issues
 - Memory management and garbage collection
 - Resource sharing and synchronization
 - High precision & accurate timing (< ms)
 - Current timer is based on nanosleep() call
 - Precise, but not accurate enough to handle (3,3) schedulers
 - Supporting of thread dispatching model:
 - RTSJ 1.1 thread dispatching model by Wellings in [3]
 - More generalized dispatching for non-priority based scheduling





Thank you for your attention

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